Box splines

Let C_N be the N dimensional hypercube:

$$C_N := \{t_1, t_2, \ldots, t_N; 0 \le t_i \le 1\}$$

We slice it with the hyperplane:

$$H_t = \{\sum_{i=1}^N t_i = t\}$$

$$Box_N(t) = volume(C_N \cap H_t).$$

Then $Box_N(t)$ is supported on $0 \le t \le N$. On each interval [k, k+1] given by a polynomial of degree N-1. But different polynomials on each interval. Their N-2 first derivatives agree at the extreme of intervals.

Remark the symmetry: $Box_N(t) = Box_N(N - t)$ We also see that

$$\frac{d}{dt}Box_N(t) = Box_{N-1}(t) - Box_{N-1}(t-1)$$

That is

$$\frac{d}{dt}Box_N(t) = (\nabla Box_{N-1})(t)$$

where ∇ is the difference operator. (We will not use this equation in this elementary talk: it holds in the distribution sense , for N > 1, because of C^{N-2} -differentiability properties)

Example: $Box_1(t)$



$Box_2(t)$



$Bot_3(t)$



Wonderful properties of Box splines

For example

$$\sum_{n\in\mathbb{Z}}Box_N(t-n)=1$$

The following pictures (see Procesi document) shows the sum of $Box_1(t-n)$, $Box_2(t-n)$, $Box_3(t-n)$ over the integers n = -1, 0, ..., 8.





This property follows right away from the geometric definition: Example: Box_2 : By drawing... Compute $\sum_n Box_2(t-n)$: We have to sum all the volumes of $x_1 + x_2 = t - n$ for any n. Now x_1 ranges between 0 and 1. Put $x_2 = t - x_1 - n$ where n is the integer $floor(t - x_1)$ so that x_2 is between 0 and 1. So it is just parametrized by $0 \le x_1 \le 1$. More generally, we will see that:

Theorem

For any polynomial P of degree strictly less than N,

$$t->\sum_{n}P(n)Box_{N}(t-n)$$

is a polynomial function of t.

N = 1

• Constant function: already seen.

• P(t) = t of degree too big:

Drawing: I think I sum over n = 0, 1, 2, 3, 4

Compute $\sum_{n} nBox_1(t-n)$ not a polynomial !!



N = 2

• Constant function: already seen.

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Function P(t) = t of degree 1: true Drawing: I think I summed over n = -1, 0, 1, 2, 3, 4, ..., 8



Series of differential operators

If $\mathbf{T}(\frac{d}{dt}) = \sum_{n=0}^{\infty} t_n (\frac{d}{dt})^n$ is a series of differential operators, we can act on polynomials.

 $(\mathbf{T}(\frac{d}{dt})P)(t)$ can be computed as for *n* large $(\frac{d}{dt})^n P = 0$. We will use the series

$$(\frac{(1 - exp(-\frac{d}{dt}))}{\frac{d}{dt}}) = 1 - \frac{1}{2}\frac{d}{dt} + \frac{1}{3!}(\frac{d}{dt})^2$$

and the inverse

Definition

We define the Todd operator

$$\mathbf{Todd}(\frac{d}{dt}) = \frac{\frac{d}{dt}}{(1 - exp(-\frac{d}{dt}))}$$
$$= 1 + \frac{1}{2}\frac{d}{dt} + \frac{1}{12}(\frac{d}{dt})^2 - \frac{1}{720}(\frac{d}{dt})^4 + \cdots$$

Integration against the Box Spline

Let *test* be a smooth function of t. Then we can integrate $test(t_1 + t_2 + \cdots + t_N)$ over the hypercube. By Fubini theorem, we obtain

$$\int_0^1 \int_0^1 test(t_1 + t_2 + \dots + t_N) d\mathbf{t} = \int_{\mathbb{R}} test(t) Box_N(t) dt$$

Definition

If test is a smooth function of t, define

• The usual convolution

$$(Box_N *_c test)(t) = \int_{u \in \mathbb{R}} test(u)Box_N(t-u)du$$

• The semi-discrete convolution

$$(Box_N *_d test)(t) = \sum_{u \in \mathbb{Z}} test(u)Box_N(t-u)$$

If test is a **polynomial**, we have (use Taylor expansion)

$$\int_0^1 test(t-u)du = (\frac{1-e^{-d/dt}}{d/dt})test(t).$$

Thus

$$(Box_N *_c test)(t) = \int_{u \in \mathbb{R}} test(u)Box_N(t-u)du$$
$$= \int_{u_1=0}^1 \cdots \int_{u_N=0}^1 test(t - (u_1 + u_2 + \cdots + u_N))du_1 \cdots du_N$$
$$= ((\frac{1 - e^{-\frac{d}{dt}}}{\frac{d}{dt}})^N test)(t).$$

Theorem

If P is a polynomial function, then

$$Box_N *_c P = ((\frac{1 - e^{-\frac{d}{dt}}}{\frac{d}{dt}})^N P)(t).$$

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If P is a polynomial function of degree < N, then

$$Box_N *_c P = Box_N *_d P == \left(\left(\frac{1 - e^{-\frac{d}{dt}}}{\frac{d}{dt}} \right)^N P \right)(t).$$

Proof

Let Box_N : it is enough to prove this for $P(t) = t^{N-1}$

$$t->\sum_n n^{N-1}Box_N(t-n)$$

is a polynomial; For smaller degree, we derivate, and we can use the recurrence formula.

I will compute

$$\sum_{n} P(n) Box_N(t-n)$$

not exactly for $P(t) = t^{N-1}$ but for the polynomial of degree N-1 given by

$$t
ightarrow rac{(t+1)(t+2)\cdots(t+N-1)}{(N-1)!}$$

Consider the standard (N-1)-dimensional simplex (dilated)

$$S_{N-1}(t) = \{t_i \ge 0; t_1 + t_2 + \dots + t_N = t.\}$$

 $S_{N-1}(t)$ has volume $\frac{t^{N-1}}{(N-1)!}$. If t = n is an integer the number of integral points in $S_{N-1}(t)$ is

$$\mathcal{P}_N(n) = \frac{(n+1)\cdots(n+N-1)}{(N-1)!}$$

Now let us integrate over the first quadrant $t_i > 0$, the function $e^{-(t_1+t_2+\cdots+t_N)y}$.

$$I=\int_{t_1>0}\cdots\int_{t_N>0}e^{-(t_1+t_2+\cdots+t_N)y}dt_1dt_2\cdots dt_N.$$

Using Fubini, we compute *I* by integrating first over the simplices $S_{N-1}(t)$ then over *t*, thus

$$I := \int_{t>0} \frac{t^{N-1}}{(N-1)!} e^{-ty} dt$$

But we can also decompose the quadrant in cubes $[n_1, n_2, n_3, \ldots, n_N] + Hypercube \ n_i = 0, 1, \ldots, \ O \le t_i \le 1$ and obtain that I is equal to

$$\sum_{\mathbf{n}} \int_{t_1=0}^1 \cdots \int_{t_N=0}^1 e^{-((n_1+t_1)-(n_2+t_2)-(n_3+t_3)-\cdots-(n_N+t_N))y} dt_1 dt_2 \cdots dt_N.$$

$$= \int_{t \in \mathbb{R}} \sum_{n_i} e^{-(\sum n_i)y} Box_N(t) e^{-ty} dt$$

$$=\int_{t\in\mathbb{R}}\sum_{n\geq 0}\mathcal{P}_{N}(n)e^{-ny}Box_{N}(t)e^{-ty}dt$$

$$=\int_{t\in\mathbb{R}}\sum_{n\geq 0}\mathcal{P}_N(n)Box_N(t-n)e^{-ty}dt$$

We obtain thus that for every y > 0

$$I := \int_{t>0} \frac{t^{N-1}}{(N-1)!} e^{-ty} dt$$

and also

$$=\int_{t\in\mathbb{R}}\sum_{n\geq 0}\mathcal{P}_N(n)Box_N(t-n)e^{-ty}dt.$$

CONCLUSION: For t > 0, we have almost everywhere

$$\sum_{n\geq 0}\mathcal{P}_N(n)Box_N(t-n)=\frac{t^{N-1}}{(N-1)!}.$$

So we have it on each interval where Box_N is continuous.

Recall that we want to compute

$$\sum_{n\in\mathbb{Z}}P(n)Box_N(t-n)$$

for the polynomial

$$t
ightarrow rac{(t+1)(t+2)\cdots(t+N-1)}{(N-1)!}$$

The sum is over the integers *n* such that $t - n \le N$, as Box_N is supported on [0, N]. thus over the integers $-(N-1), -(N-2), -(N-3), \ldots, -1, 0$. But my polynomial vanishes there, and for $n \ge 0$ coincide with $\mathcal{P}_N(n)$. Thus I obtain For $t \ge 0$

$$\sum_{n\in\mathbb{Z}}P(n)Box_N(t-n)=\frac{t^{N-1}}{(N-1)!}$$

Same calculation for t < 0, using $Box_N(t) = Box_N(N - t)$

$$\sum_{n\in\mathbb{Z}}P(n)Box_N(t-n)=\sum_{n\in\mathbb{Z}}P(n)Box_N(N+n-t)$$

$$=\sum_{n\in\mathbb{Z}}P(-m-N)Box_N(-t-m)$$

Remark that $P(-m - N) = (-1)^{(N-1)}P(m)$ and we obtain the same formula For t < 0

$$\sum_{n\in\mathbb{Z}}P(n)Box_N(t-n)=\frac{t^{N-1}}{(N-1)!}.$$

It remains to see that

$$(rac{(1-e^{-\partial_t})}{\partial_t})^N$$
 binomial $(t+N-1,N-1)=rac{t^{N-1}}{(N-1)!}.$

For example, by induction.

Using distributions, we could have seen directly that

$$(\frac{d}{dt})^N(Box_N*_d P)=0$$

so that the result is a polynomial of degree < N. as

$$\frac{d}{dt} * (Box_N *_d P) = (\nabla Box_{N-1}) *_d P = Box_{N-1} * \nabla P.$$

Let f be any function on \mathbb{Z} . Consider the function on \mathbb{R}

$$F(t) := (Box_N *_d f)(t) := \sum_n f(n)Box_N(t-n).$$

Then F is a locally polynomial function of t (on each interval it is given by a polynomial function of t). We can derivate F over any open interval by any series of differential operator $P(\frac{d}{dt})$.

The Todd operator

Theorem

Let f be any function on \mathbb{Z} : Let

$$F_N(t) = (Box_N *_d f)(t) := \sum_n f(n)Box_N(t-n).$$

Then $F_N(t)$ is a function on \mathbb{R} , polynomial on each interval. Then f(n) (n an integer) is obtained by the limit from the right of

$$f(n) = \lim_{\epsilon > 0, \epsilon - > 0} ((Todd(d/dt))^N (Box_N *_d f))(n + \epsilon)$$

For some very interesting cases, the function $(Box_N *_d f)$ is known and related to the "classical" geometry.

If we know regularity properties of $(Box_N *_d f)$, then we deduce regularities properties for f.

Going from F_N to f is going from the classical mechanics to quantum mechanics.

We will see examples later.

Proof

We want to prove this equation, for n = 0 (enough). Then $F_N(t)$ for t > 0 near 0 (HERE I USE THE LIMIT on the RIGHT) involves only the values $f(0), f(-1), f(-2), \dots f(-(N-1))$ of f, indeed $Box_N(t + N) = 0$ for t > 0, as Box_N is supported on [0, N]. We can choose a unique polynomial P of degree N - 1 which coincide with f at $0, -1, \dots, -(N-1)$. Near t > 0, small,

$$(Box_N *_d f)(t) = (Box_N *_d P)(t)$$

Differentiate with the reverse operator, we obtain our identity $f(0) = (Todd(\frac{d}{dt}))^N (Box_N *_d f)(0).$

A wonderful property of the Box spline

Apply this to f(n) = 0 except for n = 0 where f(0) = 1. THAT IS

 $f = \delta_0$.

Then

Theorem Consider the locally polynomial function

 $((Todd(d/dt))^N Box_N)(t)$

Then

$$((Todd(d/dt))^N Box_N)|_{\mathbb{Z}} = \delta_0$$

(limits from the right)

Examples



We have used essentially the relation that $\mathcal{P}_N(n)$, the number of integral points in the the standard simplex can be obtained from the volume $vol(S_N(t))$ by applying the Todd operator. Consider $M := P_{N-1}(\mathbb{C})$ realized by

$$\{(z_1, z_2, \ldots, z_N); \sum_i |z_i|^2 = t\}/e^{i\theta}$$

with symplectic form $\Omega_t = tc$. Here *c* is the Fubini-Study canonical 2-form on *M*, with $\int c^{N-1} = 1$. We compute

$$vol(M_t) := \int_M e^{tc} = \frac{t^{N-1}}{(N-1)!}$$

Let t = n an integral value,

Let \mathcal{L}_n be the line bundle on M with holomorphic sections polynomials on degree n: Thus $H^0(M, \mathcal{L}_n)$ has basis $z_1^{n_1} \dots z_N^{n_N}$ with $n_i \ge 0$; $\sum n_i = n$. That is the number of integral points in $S_{N-1}(n)$.

If we apply the Todd operator

$$Todd(rac{d}{dt})\int_{M}e^{tc}$$

we obtain

$$\int_M e^{tc} (\frac{c}{1-e^{-c}})^N$$

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For t = n, we then obtain

$$Todd(rac{d}{dt})\int_{M}e^{tc}|_{t=n}=\int_{M}chern(L_{n})Todd(M).$$

The Riemann Roch theorem asserts that this is the dimension of $H^0(M, L_n)$ (no higher cohomology).

The Mother Formula

CONCLUSION: The relation

$$((\mathit{Todd}(rac{d}{dt}))^N\mathit{Box}_N)|_{\mathbb{Z}} = \delta_0$$

is the mother formula:

Children

•: Inversion formula for semi-discrete convolution

• Riemann-Roch theorem for $P_{N-1}(\mathbb{C})$.

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Multiplicities formulae: last talk.